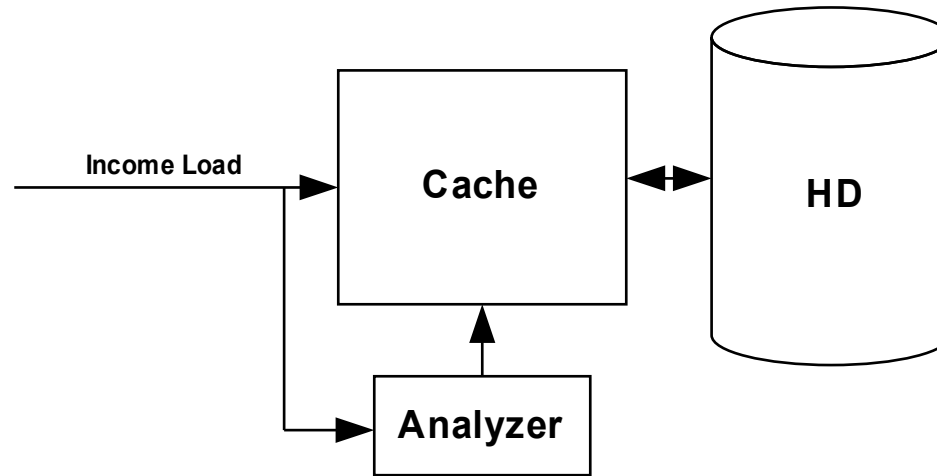


Fast Method for Dependency Mining

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Use Case



- Analyzer detects dependent blocks by observing income load
- If some block was referenced cache prefetch depended blocks found by analyzer

Dependency of Blocks

- Let $\text{support}(ab, \delta)$ – the number of times when after reference to block a reference to block b is occurred with difference at most δ references
- Let $\text{support}(a)$ – the number of reference to block a
- Block b is called dependent from block a if
 - $\text{support}(a) \geq m$
 - $p(b | a) \geq p_{\text{threshold}}$, where $p(b | a) = \frac{\text{support}(ab, \delta)}{\text{support}(a)}$
- Proposed definition of dependency between pair of blocks could be easily generalized on the dependency between sequence of blocks, e.g. a_1, a_2, \dots, a_{l-1} and some block, e.g. a_l .
e.g. al. $\underbrace{a \ c \ b}_{\delta=1} \ e \ f \ g \ h \ \underbrace{a \ b \ k}_{\delta=1} \ l \ p \ \underbrace{a \ r \ b}_{\delta=1} \ e$

Basic Method (C-Miner)

Frequent Sequence Search

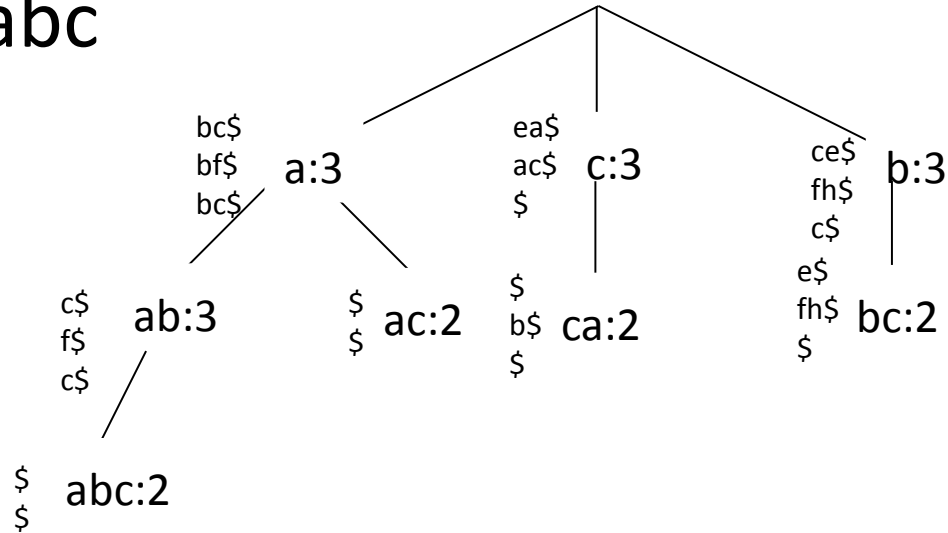
abceabfhcabc

$\delta = 1, m=2$

a: bc\$, bf\$, bc\$

b: ce\$, fh\$, c\$

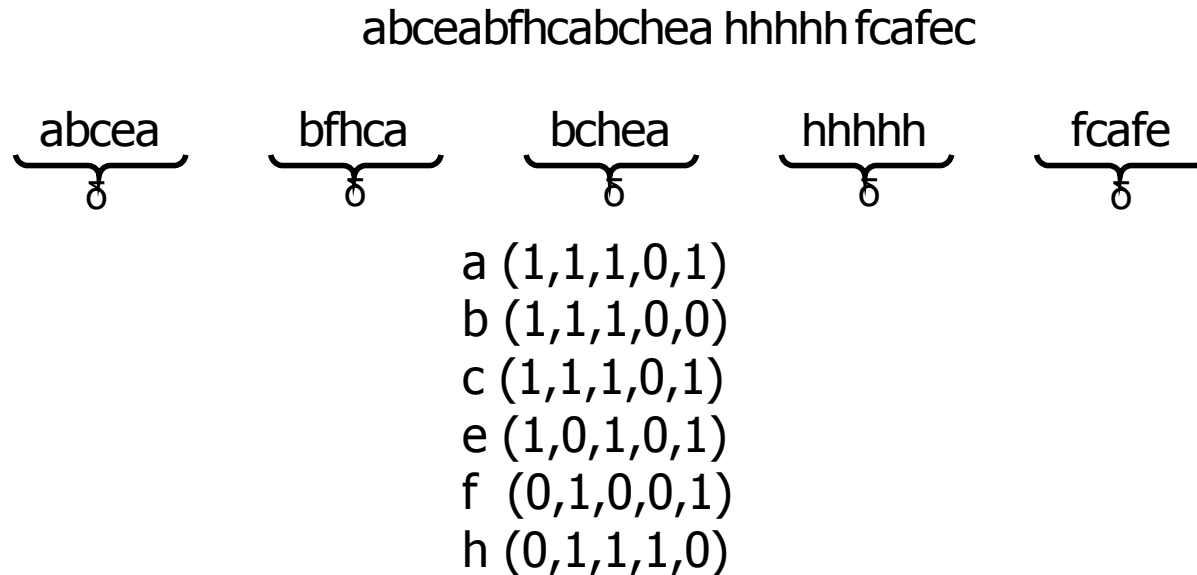
c: ea\$, ab\$, \$



Frequent Subsequences = {ab, abc, ac, ca, a, b, c}

- Number of children in each node could reach number of unique reference in a sequence (n). Then complexity of algorithm is $O(n\delta)$
- Size of stored data structure could exceed the sequence length

Proposed Approach



- Source sequence is split on the subsequences of equal length δ
- Each unique block is assigned activity vector
 - Length – the number of subsequences of length δ

Expected Results

- Implementation of LRU cache simulator
- Implementation of C-Miner algorithm
- Implementation of proposed approach
- Estimation of rules quality obtained by proposed approach
- Performance evaluation of LRU with rules obtained by
 - C-Miner
 - Proposed approach

Q&A